

Year	1990	1991	1992	1993	1994	1995	1996	1997	1998	1999	2000	2001	2002	2003	2004	2005	2006	2007	2008	2009	2010	2011	2012	2013	2014	2015	2016	2017	2018	2019	2020	2021	2022	2023	2024	2025	2026	2027	2028	2029	2030	2031	2032	2033	2034	2035	2036	2037	2038	2039	2040	2041	2042	2043	2044	2045	2046	2047	2048	2049	2050	2051	2052	2053	2054	2055	2056	2057	2058	2059	2060	2061	2062	2063	2064	2065	2066	2067	2068	2069	2070	2071	2072	2073	2074	2075	2076	2077	2078	2079	2080	2081	2082	2083	2084	2085	2086	2087	2088	2089	2090	2091	2092	2093	2094	2095	2096	2097	2098	2099	2100
1990	1991	1992	1993	1994	1995	1996	1997	1998	1999	2000	2001	2002	2003	2004	2005	2006	2007	2008	2009	2010	2011	2012	2013	2014	2015	2016	2017	2018	2019	2020	2021	2022	2023	2024	2025	2026	2027	2028	2029	2030	2031	2032	2033	2034	2035	2036	2037	2038	2039	2040	2041	2042	2043	2044	2045	2046	2047	2048	2049	2050	2051	2052	2053	2054	2055	2056	2057	2058	2059	2060	2061	2062	2063	2064	2065	2066	2067	2068	2069	2070	2071	2072	2073	2074	2075	2076	2077	2078	2079	2080	2081	2082	2083	2084	2085	2086	2087	2088	2089	2090	2091	2092	2093	2094	2095	2096	2097	2098	2099	2100	

BACKGROUND OF THE INVENTION

The present invention relates in general to a network management system. More particularly, the invention relates to a network management system for
5 managing a network system to which a plurality of networks the address systems of which are different from one another are connected.

IP addresses as network addresses which are currently most widely utilized in the Internet and the
10 like are prescribed in accordance with the protocol called the IPv4 (Internet Protocol ver. 4) and also are addresses each having the address space of 32 bits. The network addresses must be uniquely allocated to apparatuses connected to the network, respectively, and
15 hence the allocation of the network addresses to the apparatuses connected to the Internet has been carried out by the organ called NIC or the like in such a way as to become unique. Since in the IPv4, each of the IP addresses has the address space of 32 bits, the
20 addresses of 2 to the 32-th power, i.e., four billions at maximum can be theoretically allocated. However, as the number of apparatuses connected to the Internet is increased, it becomes difficult to allocate the IP addresses of the IPv4 to all of the apparatuses
25 connected to the Internet.

The technique which is widely used as the technique for solving the above-mentioned problem is the method wherein the address translation function described in RFC1631 (The IP Network Address Translator) decided by the IETF is combined with the private network described in RFC1918 (Address Allocation for Private Internets). The method is established on the assumption that for example, all of the apparatuses which are present in the local network such as the network within industry are not necessarily connected to the external network such as the Internet. In other words, first of all, the local network such as network within industry allocates the IP address using the address which falls within the private address described in RFC1918 to construct the network. At this time, with this private address, the associated apparatuses are not connected to the external network such as the Internet. With respect to the apparatuses connected to the external network such as the Internet, the address translator described in RFC1613 is disposed between the local network and the external network such as the Internet, and the address of the transmission/reception packet is translated through the address translation from the private address into the global address with which the access is given to the Internet, thereby making the connection to the Internet possible. Now, by the global address is meant the address which is allocated by the NIC or the like. At this time,

there is also made the device in which the infinite global addresses are effectively utilized in such a way that the global addresses are dynamically allocated to make it possible that a plurality of local nodes hold
5 one global address in common on the basis of the time-sharing.

Now, in the address translation described in RFC1631, the rewriting of the transmission source IP address and the transmission destination IP address
10 which are contained in the header of the IP packet, and the change of the check sum of the IP header which is generated along with the rewriting are recalculated to replace them with each other. As a result, the communication by TCP/IP becomes possible, and the
15 communication according to the protocol of the upper layer with respect thereto also becomes possible.

However, though in the network management protocol such as the SNMP, the IP address is contained in a protocol data unit (PDU) as well in which the data
20 is exchanged in accordance with the management protocol, this part is not subjected to the address translation in the address translator described in RFC1631.

On the other hand, in JP-A-11-187058, in
25 addition to the function of the address translation described in RFC1631, the address translator having the function of carrying out the address translation with respect to the protocol data unit as well of the

management protocol is described.

SUMMARY OF THE INVENTION

However, if the operation based on the private address and the global address is carried out by the address translator, then this operation will be carried out through the communication in accordance with the management protocol using the global address which can be used in the network on the side of the manager node. For this reason, the communication in accordance with the management protocol can be carried out only with the node which is located on the managed side and to which the global address is already allocated. However, if in the network management, the management can be carried out for the node as well which has not such a global address allocated thereto, it becomes possible to carry out the more effective management. This respect is not disclosed in the method described in JP-A-11-187058.

In the light of the foregoing, the present invention has been made in order to solve the above-mentioned problems associated with the prior art, and it is therefore an object of the present invention to provide a network management system in which the data communication in accordance with a management protocol can be carried out even between nodes each not having a global address in a network system in which networks of different address systems are connected to one another

through an address translator.

As one means for attaining the above-mentioned object, according to the present invention, there is provided a network management system for
5 managing a network system in which a first network and a second network which are different in address system from each other are connected to each other through an address translator, wherein each of the first and second networks includes a node and a management
10 protocol proxy, and each of the management protocol proxies includes a management protocol proxy data generating unit for treating, as the management protocol proxy data, a transmission source address, a transmission destination address and data in a protocol
15 data unit which are contained in a packet of a management protocol sent from a node, and an address translation unit for translating address information within a protocol data unit contained in management protocol proxy data sent from another management protocol proxy.

20 As a result, it is possible to carry out the data communication between the proxies of the management protocol, and it is also possible to carry out the data communication between the nodes each not having the global address.

25 BRIEF DESCRIPTION OF THE DRAWINGS

The above and other objects as well as advantages of the present invention will become clear

by the following description of the preferred embodiments of the present invention with reference to the accompanying drawings, wherein:

Fig. 1 is a diagram showing a basic configuration of a network system;

Fig. 2 is a diagram showing a structure of a packet of an SNMP message;

Fig. 3 is a diagram showing a structure of a packet of management protocol proxy data;

Fig. 4 is a diagram useful in explaining a function of an NAT;

Fig. 5 is a diagram showing a structure of a packet of an SNMP message;

Fig. 6 is a diagram showing a structure of a packet of an SNMP message;

Fig. 7 is a diagram showing one application example in a virtual network management system;

Fig. 8 is a diagram showing one application example in a virtual network management system;

Fig. 9 is a diagram showing one application example in a virtual network management system;

Fig. 10 is a diagram showing one application example in a virtual network management system;

Fig. 11 is a diagram showing one application example in a virtual network management system;

Fig. 12 is a block diagram showing a configuration of a management protocol proxy on the managed side;

Fig. 13 is a block diagram showing a configuration of a management protocol proxy on the managing side;

Fig. 14 is a diagram showing one example of
5 management protocol proxy data;

Fig. 15 is a block diagram showing a configuration of a proxy data disassembly/assembly unit;

Fig. 16 is a diagram showing a definition example of the other party proxy definition;

10 Fig. 17 is a flow chart useful in explaining the processing of assembling proxy data;

Fig. 18 is a flow chart useful in explaining the processing of disassembling proxy data;

Fig. 19 is a block diagram showing a configuration of an address translation processing unit;
15

Fig. 20 is a block diagram showing a configuration of an object identifier address translation unit;

Fig. 21 is a flow chart useful in explaining
20 the processing executed in a PDU analysis/translation unit;

Fig. 22 is a diagram useful in explaining the relationship among a kind of PDU, the direction of transmitting a PDU and the translation direction of the
25 address translation;

Fig. 23 is a flow chart useful in explaining the processing in a translation subject object identifier extraction unit;

Fig. 24 is a diagram showing a definition example of the address translation definition;

Fig. 25 is a diagram showing an example of a configuration of a virtual network management system;

5 Fig. 26 is a block diagram showing a configuration of a management protocol proxy;

Fig. 27 is a block diagram showing a configuration of a management protocol proxy;

Fig. 28 is a diagram showing an example of a
10 configuration of a virtual network management system;
and

Fig. 29 is a block diagram showing a configuration of a management protocol proxy on the managed side.

15 DESCRIPTION OF THE EMBODIMENTS

The preferred embodiments when an SNMP is applied to a management protocol will hereinafter be described in detail with reference to the accompanying drawings.

20 Fig. 1 is a diagram showing the concept of a network management system which will be described in the present embodiment.

The present system includes a global network 10, a private network A 30a and a private network B
25 30b. Then, the private network A 30a is connected to the global network 10 through an NAT 20a for carrying out the address translation, while the private network

B 30b is connected to the global network 10 through an NAT 20b for carrying out the address translation.

A node 80 for carrying out the monitoring is connected to the private network A 30a. This node 80
5 includes a manager 40 for executing the processing for monitoring managed nodes, and a management protocol proxy 60a for carrying out the generation of proxy data of the management protocol, the address translation within a protocol data unit contained in the proxy
10 data, and the like.

In addition, nodes 50 and 90 as the managed nodes, and the management protocol proxy server 60b for carrying out the generation of proxy data of the management protocol, the address translation within a
15 protocol data unit contained in the proxy data, and the like are connected to the private network B 30b.

In such a configuration, in the present network management system, the node 80 for carrying out the monitoring is adapted to manage the nodes 50 and 90
20 as the managed nodes.

When the data communication is intended to be carried out with the node 50 as the managed node through the manager 40 of the node 80 of the private network 30a, the data communication is carried out
25 between the management protocol proxy 60a of the node 80 and the management protocol proxy server 60b of the private network B 30b. This becomes equivalent to that a virtual communication path called an interproxy

communication path 70 is formed.

If such a configuration is adopted, then the data communication can be carried out between the networks when the address used in the private network A 30a is different from that used in the private network B 30b, i.e., when their address systems are different from each other. In addition, even in the case where a fire wall or the like which is adopted to block the passage of any of the addresses of the private networks and the address of the global network is formed between the private network A 30a and the private network B 30b, the data communication can be made possible.

The operation of the overall network system shown in Fig. 1 will hereinbelow be described more concretely.

Fig. 2 is a diagram showing an SNMP message packet of a management protocol SNMP of the node. Fig. 3 is a diagram showing a packet of management protocol proxy data of the management protocol proxy server. That is, the management protocol proxy 60a and the management protocol proxy server 60b store the transmission source information and the transmission destination information which have been respectively held as the transmission destination and the transmission source within the header of an IP layer as a network layer in the management protocol proxy data corresponding to the data of an application layer from the packets of the management protocol SNMPs which have

been respectively sent from the manager 40 of the node 80, and the node 50 and the node 90 as the managed nodes. Further, as for a transport layer of the packet of the management protocol proxy data, for example, a TCP of a connection type is employed. In addition, the transmission destination and the transmission source of the management protocol proxy data itself become the management protocol proxy or the management protocol proxy server.

Now, the description will hereinbelow be given with respect to the flow of the data when the communication is carried out from the manager 40 to the node 50 or the node 90 as the managed node and in the direction opposite thereto, i.e., from the node 50 or the node 90 as the managed node to the manager 40 through the management protocol proxy 60a and the management protocol proxy server 60b. At the time when the data of the management protocol sent by the node 40 as the manager has been delivered to the management protocol proxy 60a, the management protocol proxy 60a combines the data itself of the management protocol with the information exhibiting the essential transmission destination and transmission source of the data of the management protocol to generate the management protocol proxy data and sends the management protocol proxy data thus generated to the management protocol proxy server 60b. Then, the management protocol proxy server 60b fetches the data of the management protocol

and the information exhibiting the essential transmission destination from the received management protocol proxy data to reconstruct the packet of the management protocol to send the packet thus

5 reconstructed to the node 50 as the managed node which is the essential transmission destination. At this time, the transmission source of the packet of the management protocol is made the management protocol proxy server 60b, whereby the response from the node 50
10 is sent to the management protocol proxy server 60b. Then, the management protocol proxy server 60b which has received the response packet of the management protocol from the node 50 translates that response packet into the management proxy protocol data to
15 return the resultant management proxy protocol data back to the management protocol proxy 60a. Then, the management protocol proxy 60a reconstructs the response packet of the management protocol from the management protocol proxy data to return the resultant response
20 packet back to the manager 40.

In such a manner as described above, in the environment in which the communication can not be directly carried out between the nodes using the management protocol, the communication in accordance
25 with the management protocol can be carried out through the management protocol proxy and the management protocol proxy server.

Further, each of the management protocol

proxy and the management protocol proxy server includes an address translation function of translating the address of the PDU part of the management protocol. As a result, the address of the node contained in the
5 protocol data unit (hereinafter, referred to as "the PDU" for short, when applicable) of the management protocol of the node is translated into the virtual address which is used to manage the network, whereby the management can be carried out as if the node has
10 the virtual address. The virtual address which is used only in the network management at this time is referred to as "the management address" for short when applicable in the present embodiment for the sake of convenience.

15 By the way, while in the configuration shown in Fig. 1, the description has been given with respect to the case where the management protocol proxy 60a is realized in the form of the program, similarly, the function of the management protocol proxy server 60b
20 can also be realized in the form of the program. In this case, the program is recorded in a magnetic disk, an optical disc or a magneto-optical disc from which the data can be read out by a computer, and the node which takes charge of the function of the management
25 protocol proxy server is adapted to read out the program to execute the program.

Next, the NAT shown in Fig. 1 will hereinbelow be described. Each of the NATs shown in Fig. 1

is in conformity to RFC1631.

Fig. 4 is a diagram useful in explaining the function of the NAT. As shown in the figure, the description will hereinbelow be given with respect to

5 the case where in the network in which a global network 10 and a private network 30 are connected to each other through an NAT 20, an address G0 is allocated as a global address 200 to a node 40, an address L1 is allocated as a private address 210 to the managed node,

10 and an address G1 is allocated as the global address 200 to the managed node. In this case, a translation table 230 in which the global address and the private address are made correspond to each other is provided in the NAT 20. As a result, the packet, with respect

15 to the address G1, is transmitted from the node 40 side to the node 50 side. The NAT 20 translates the transmission destination address of the packet directed to the address G1 from the global address G1 to the private address L1 in accordance with the translation

20 table 230 to transmit the resultant address to the private network side. That is, when having reached the NAT 20 from the node 40 side as shown in Fig. 5, the packet of interest is the packet in which the information of the address G1 for the transmission destination

25 is contained as the header information of the IP layer part corresponding to the network layer, while when being relayed from the NAT 20 towards the node on the private network side, is transmitted as the packet in

which the information of the address L1 is contained as the header information of the IP layer part as shown in Fig. 6 in the transmission destination. On the other hand, in the case where the packet is transmitted from the node 50 on the private network side towards the node 40 on the global network side, when having reached the NAT 20 from the node 50, the packet of interest is the packet in which the information of L1 for the transmission source is contained as the header information of the IP layer part, while when being relayed from the NAT 20 towards the node 40 on the global network side, is transmitted as the packet in which the information of G1 for the transmission source is contained as the header information of the IP layer part.

On the basis of such a function of the NAT 20, as for the setting of the network of the node 50 itself, the setting has only to be made in such a way that the network of the node 50 itself has the private address L1. Then, when the node 50 is intended to communicate with another apparatus within the private network, the communication can be carried out using the private address L1. In addition, when the communication is intended to be carried out with the apparatus on the global network side, the communication can be carried out using the global address G1 allocated by the NAT 20.

By the way, in this case, the NAT has been

described as the translator having the function of carrying out the address translation of the IP layer (the address translation of the PDU part of the management protocol can not be carried out). In the network utilizing such an NAT, as has already been described, the address of the PDU part of the management protocol is translated in the management protocol proxy server.

The more concrete network system will herein-
below be described.

Fig. 7 is one of application examples in the network management system and shows a configuration in the case where the management is carried out by using the proper global address allocated to a managed side node as the address which is used to manage the network.

A managed node a 50a has the address L1 as the private address 210. This address of L1 is the private address, and hence is the address which can be used only in the private network B 30b.

Furthermore, as for address translation 230 in the NAT 20b, the global address G1 is made correspond to the private address L1, and the address G1 is statically allocated as the global address to the managed node a 50a. When the global network 10 or the private network A 30a is intended to communicate directly with the managed node a 50a, the communication is carried out using the global address G1.

Now, as for address translation 220 in the management protocol proxy 60b, the management address G1 is made correspond to the private address L1, whereby with respect to the management protocol data which is to be sent from the managing side to the managed side, the global address G1 is translated into the private address L1, while with respect to the management protocol data which is to be sent from the managed side to the managing side, the private address L1 is translated into the global address G1.

As a result, if reference is made to the data of the management protocol using the management protocol from the manager side, the managed node a 50a seems to be as if it has the global address G1.

For this reason, as for the management information 240 in the manager 40, the managed node a 50a is made correspond to the apparatus having the address G1, and hence the network management can be carried out using the address G1.

Next, Fig. 8 shows the case where the global address to be allocated to the managed node by the NAT 20b is dynamically allocated thereto.

In general, the number of global addresses which can be used for the external connection in the private networks which are connected through the NATs is less than the number of apparatuses within the private networks. As for the method of utilizing effectively the less global addresses, there is

employed the method wherein the allocation of the global address to the apparatus is carried out only for a period of time when the apparatus of interest is connected to the outside, and at the time when the connection of the apparatus of interest to the outside has been completed, the global address which the apparatus of interest has used is adapted to be reutilized by another apparatus. In the case where the global addresses are dynamically allocated in such a manner, in general, the global address which is allocated to a certain one apparatus differs as the case may be.

In such a case, as the address translation 230 in the NAT 20b, the global address Gx is made correspond to the private address L1. Then, the global address Gx is the address which is selected among the addresses of the fixed choices as the case may be, and hence the address Gx is dynamically allocated as the global address to the managed node a 50a at least at that time. What address is allocated is determined by the NAT 20b.

Now, if the address Gx is used as the management address as it is, then the address which is made correspond as the management information 240 in the manager 40 to the managed node a 50a changes as the case may be, and hence the continuation of the management can not be maintained so that the proper network management can not be carried out. This is a problem.

Then, as for the address translation 220 in the management protocol proxy 60b, as shown in Fig. 8, the virtual address of V1 which is completely different from the address Gx is statically allocated. As a
5 result, as for the management information 240 in the manager, the managed node a 50a is made correspond to the apparatus having the address V1 and hence the network management can be properly carried out.

Fig. 9 shows an example in the case where
10 there are a plurality of private networks on the managed side, and the private addresses of the managed nodes within these private networks compete with one another.

A managed node b 50b is present in the
15 private network B 30b and has the private address L1. On the other hand, a managed node c 50c is present in a private network C 30c and has the private address L1. While the respective addresses conflict with each other, since the private address is used only within
20 the associated one of the private networks in terms of the communication, the networks are not confused at all.

However, if the management protocol data is directly acquired from the managed node b 50b and the
25 managed node C 50c using the management protocol, since both of these apparatuses respond to the information as the apparatus having the private address L1, there arises the problem that the manager is confused so that

the network management can not be properly carried out.

Then, as shown in Fig. 9, with respect to the managed node b 50b, the management address V1 is made correspond to the private address L1 during the address translation 220b in the management protocol proxy 60b, while with respect to the managed node C 50c, the management address V2 is made correspond to the private address L1 during the address translation 220c in the management protocol proxy 60c. That is, the management address V1 and the management address V2 are statically allocated to the managed node b 50b and the managed node c 50c, respectively.

As a result, as for the management information 240 in the manager, the managed node b is decided as the apparatus having the management address V1, and the managed node c is decided as the apparatus having the management address V2 so that the network management can be properly carried out.

Fig. 10 shows an example in the case where the management address is allocated to even the apparatus to which the global address is not allocated, i.e., which is not connected to the outside in order to carry out the network management in the private network.

A managed node a 50a has the address L1 as the private address 210a. In addition, as for the address translation 230 in the NAT 20b, the global address G1 is made correspond to the private address

L1, and the address G1 is statically allocated as the global address to the managed node a 50a. When the global network 10 or the private network A 30a intends to communicate directly with the managed node a 50a,
5 the communication is carried out using the global address G1.

On the other hand, the managed node b 50b has the address L1 as the private address 210b. However, as for the address translation 230 in the NAT 20b, the
10 global address corresponding to the private address L2 is not defined, and hence the global network 10 or the private network A 30a can not communicate directly with the managed node b 50b. But, in this case as well, the manager 40 can exchange the information with the
15 managed node 50 as well in accordance with the management protocol through the management protocol proxy.

At this time, during the address translation 220 in the management protocol proxy 60b, the management address V1 is made correspond to the private
20 address L1, and also the management address V2 is made correspond to the private address 22. That is, the management address V1 is statically allocated to the managed node a 50a, and also the management address V2 is statically allocated to the managed node b 50b. As
25 a result, as for the management information 240 in the manager, the managed node a 50a is decided as the apparatus having the management address V1, and the managed node b 50b is decided as the apparatus having

the management address V2 in order to carry out the network management.

Fig. 11 shows an example in which when the management address is allocated to even the apparatus to which the global address is not allocated, i.e., which is not connected to the outside in the private network in order to carry out the network management, with respect to the apparatus to which the global address is allocated, the global address is used for the management address, while with respect to the apparatus to which the global address is not allocated, the virtual address is used therefor.

The managed node a 50a has the address L1 as the private address 210a. In addition, as for the address translation 230 in the NAT 20b, the global address G1 is made correspond to the private address L1, and the address G1 is statically allocated as the global address to the managed node a 50a. When the global network 10 or the private network A 30a intends to communicate directly with the managed node a 50a, the communication is carried out using the global address G1.

On the other hand, the managed node b 50b has the address L2 as the private address 210b. However, as for the address translation 230 in the NAT 20b, the global address corresponding to the private address L2 is not defined, and hence the global network 10 or the private network A 30a can not communicate directly with

the managed node b 50b. But, in this case as well, the manager 40 can exchange the information with the managed node 50b as well in accordance with the management protocol through the management protocol proxy.

5 At this time, during the address translation 220 in the management protocol proxy 60b, the management address G1 is made correspond to the private address L1, and also the management address V2 is made correspond to the private address L2. That is, the
10 management address G1 is statically allocated to the managed node a 50a and also the management address V2 is statically allocated to the managed node b 50b. As a result, as for the management information 240 in the manager, the managed node a 50a is decided as the
15 apparatus having the address G1, and the managed node b 50b is decided as the apparatus having the address V2 in order to carry out the network management.

Subsequently, the configuration of the management protocol proxy will hereinbelow be described
20 with reference to Fig. 12 and Fig. 13.

Fig. 12 is a block diagram showing a configuration of the management protocol proxy on the managed side.

The management protocol proxy 60b on the
25 managed side includes: an interproxy communication unit 61 for processing the establishment of the interproxy communication path 70 between the management protocol proxy 60a on the managing side and the unit 61 and the

transmission/reception of the management protocol proxy
data; a proxy data assembly/disassembly unit 62 for
processing the disassembly/assembly management protocol
proxy data; an address translation processing unit 63
5 for subjecting the address information within the PDU
of the management protocol to the address translation;
an ASN.1MIB define statement and an address translation
definition 65 each of which becomes an input to the
address translation processing unit; and an SNMR
10 message transmission/reception unit 66 for
transmitting/receiving the SNMP message between the
managed node and the unit 66.

At the time when the message of the manage-
ment protocol issued from the manager 40 has been
15 translated into the management protocol proxy data by
the management protocol proxy 60a on the managing side
to be transmitted to the management protocol proxy on
the managed side, first of all, the interproxy
communication unit 61 receives the management protocol
20 proxy data transmitted thereto to deliver the manage-
ment protocol proxy data thus received to the proxy
data assembly/disassembly unit 62. Then, the proxy
data assembly/disassembly unit 62 disassembles the
received management protocol proxy data to deliver the
25 management protocol proxy data thus disassembled to the
address translation processing unit 63. Then, the
address translation processing unit 63 subjects the
transmission source address data and the transmission

destination address data of the management protocol proxy data thus delivered thereto and the address information within the PDU of the management protocol into the address translation in accordance with the

5 ASN.1MIB define statement 64 and the address translation definition 65 to deliver the translation result to the proxy data disassembly/assembly unit 62. Then, the proxy data assembly/disassembly unit 62 fetches the transmission destination information, the transmission
10 source information, and the PDU of the management protocol from the proxy data to deliver the information and the PDU thus fetched to the SNMP message transmission/reception unit 66. Then, the SNMP message transmission/reception unit 66 transmits the PDU of the
15 management protocol to the managed node 50 which has been specified as the transmission destination. That is, the SNMP message transmission/reception unit 66 transmits the SNMP message to the managed node 50. Then, the managed node 50 returns the response
20 corresponding to the SNMP message thus transmitted thereto back to the SNMP message transmission/reception unit 66. Then, the SNMP message transmission/reception unit 66 delivers the SNMP message of the response thus received, and the information of the transmission
25 source and transmission destination thereof to the proxy data assembly/disassembly unit 62. Then, the proxy data assembly/disassembly unit 62 delivers the information of the transmission destination and the

transmission source, and the PDU of the management
protocol as the response to the address translation
processing unit 63. Then, the address translation
processing unit 63 subjects the transmission source
5 address data and the transmission destination address
data of the management protocol proxy data delivered
thereto, and the address information within the PDU of
the management protocol to the address translation to
deliver the translation result to the proxy data
10 disassembly/assembly unit 62. Then, the proxy data
disassembly/assembly unit 62 assembles the information
of the transmission destination and the transmission
source, and the PDU of the management protocol in the
form of the management protocol proxy data to deliver
15 the data thus assembled to the interproxy communication
unit 61. Then, the interproxy communication unit 61
transmits the management protocol proxy data to the
management protocol proxy 60a on the managing side, and
also the management protocol proxy 60a on the managing
20 side returns the PDU of the management protocol back to
the manager 40.

Now, the ASN.1MIB define statement 64 is the
MIB define statement described by the ASN.1 (Abstract
Syntax Notation One) which is the standard description
25 method for the MIB object described in RFC1212 (Concise
MIB Definition) and the like. In general, the MIB
define statement by ASN.1 is widely open to the public
by a person who had defined the MIB module thereof. In

the present embodiment, since the address contained in Variable-Bindings of the protocol data unit of the management protocol is translated using the information which is obtained by analyzing the MIB define statement by ASN.1, the special define statement becomes unnecessary so that the configuration becomes simpler.

Fig. 13 is a block diagram showing a configuration of the management protocol proxy on the managing side.

10 The management protocol proxy 60a on the managing side includes: an SNMP message transmission/reception unit 66 for transmitting/receiving the SNMP message to/from an SNMP manager 41 on the manager 40; a proxy data assembly/disassembly unit 62 for processing
15 the disassembly/assembly of the management protocol proxy data; and an interproxy communication unit 61 for processing the establishment of the interproxy communication path 70 distributed between the management protocol proxy 60b on the managed side and the unit 61,
20 and the transmission/reception of the management protocol proxy data.

At the time when the SNMP message issued from the SNMP manager 41 on the manager 40 has been delivered to the management protocol proxy 60a on the
25 managing side, the SNMP message transmission/reception unit 60 receives the SNMP message to deliver the SNMP message thus received to the proxy data assembly/disassembly unit 62. Then, the proxy data assembly/

disassembly unit 62 assembles the management protocol proxy data from the SNMP message thus delivered thereto, and the information of the transmission destination and the transmission source thereof to

5 deliver the management protocol proxy data thus assembled to the interproxy communication unit 61.

Then, the interproxy communication unit 61 transmits the management protocol proxy data thus delivered thereto to the management protocol proxy 60b on the

10 managed side. Further, the interproxy communication unit 61 receives the management protocol proxy data of the response returned from the management protocol proxy 60b on the managed side to deliver the management protocol proxy data thus received to the proxy data

15 assembly/disassembly unit 62. Then, the proxy data assembly/disassembly unit 62 fetches the information of the transmission source and the transmission destination, and the SNMP message from the management protocol proxy data to deliver the information and the message

20 thus fetched to the SNMP message transmission/reception unit 66. Then, the SNMP message transmission/reception unit 66 returns the SNMP message back to the SNMP message 41 on the manager 40 in accordance with the delivered information.

25 Fig. 14 shows one example of the management protocol proxy data which is transmitted/received on the interproxy communication path 70 distributed between the management protocol proxy 60a on the

manager side and the management protocol proxy 60b on the managed side, and also shows an example in the case where the management protocol is the SNMP. In this case, the management protocol proxy data is the data
5 containing therein the information of the transmission source of the SNMP message, the information of the transmission destination of the SNMP message, and the SNMP PDU.

Fig. 2 is a diagram showing the packet of the
10 normal SNMP message. By the way, the packet is shown only with respect to the upper part with respect to the IP layer as the network layer. In the normal SNMP packet, the information of the transmission source and the information of the transmission destination at the
15 IP layer level become the transmission source and the transmission destination of the SNMP message itself as they are.

Fig. 3 is a diagram showing the packet of the management protocol proxy data in the case where the
20 management protocol is the SNMP. By the way, the packet is shown only with respect to the upper part with respect to the IP layer as the network layer. In the packet of the management protocol proxy data, the transmission source or the transmission destination at
25 the IP layer level is one of the management protocol proxies which are present in the both sides of the interproxy communication path 70, respectively, and the data of the transmission source and the transmission

destination of the SNMP message is contained in the form of the management protocol proxy data corresponding to the application layer in the packet. Therefore, the virtual address which is used for the address translation is not used as the transmission destination or the transmission source address of the actual communication packet. For this reason, even if the virtual address is used which is not allocated from the organ such as the NIC to the organization thereof, this does not impede the communication at the IP layer, i.e., at the network layer level at all.

In such a manner as described above, the addresses of the SNMP message transmission source and transmission destination on the management protocol proxy data are subjected to the address translation, whereby the virtual address which is not the proper global address can be used in the management address, and hence the private network can be subjected to the network management including the apparatuses each not having the global address.

Next, the proxy data assembly/disassembly unit 62 will hereinbelow be described with reference to Fig. 15, Fig. 17 and Fig. 18.

Fig. 15 is a block diagram showing a configuration of the proxy data assembly/disassembly unit 62.

The proxy data assembly/disassembly unit 62 includes: an assembly/disassembly processing unit 68

for executing the assembly/disassembly processing; and
the other party proxy definition 69 in which the
correspondence relationship between the transmission
destination of the SNMP message and the other party to
5 which the SNMP message is to be transmitted is defined.

Fig. 16 shows a definition example of the
other party proxy definition.

A definition line 311 is a definition line
exhibiting that the SNMP message in which a first octet
10 of the transmission destination address is 100 is
transmitted to the management protocol proxy in which
the address is 200.10.20.30.

A definition line 312 is a definition line
exhibiting that the SNMP message in which a first octet
15 of the transmission destination address is 101 and a
second octet thereof is 10 is transmitted to the
management protocol proxy in which the address is
200.10.20.30.

A definition line 313 is a definition line
20 exhibiting that the SNMP message in which a first octet
of the transmission destination address is 10, a second
octet thereof is 20, and a third octet thereof is 80 is
transmitted to the management protocol proxy in which
the address is 230.51.62.72.

25 A definition line 314 is a definition line
exhibiting that the SNMP message in which the trans-
mission destination address is 120.60.11.8 is
transmitted to the management protocol proxy in which

the address is 230.51.62.72.

By the way, the transmission address at this time is represented using the management address.

Fig. 17 shows a flow chart useful in explaining the proxy data assembly processing.

In Step 151, the SNMP message is received from the SNMP message transmission/reception unit.

In Step 152, both of the transmission source address and the transmission destination address are fetched from the IP header part of the SNMP message.

In Step 153, the SNMP PDU is fetched from the SNMP message.

In Step 154, the transmission source address, the transmission destination address and the SNMP PDU which have been fetched from the IP header part of the SNMP message are stored in the management protocol proxy data. In Step 155, the other party proxy address which is made correspond to the transmission destination address fetched from the IP header part of the SNMP message is retrieved in the other party proxy definition, and the other party proxy address of interest is decided as the transmission destination of the protocol proxy data.

In a manner as described above, the proxy data assembly/disassembly unit assembles the management protocol proxy data.

Fig. 18 shows a flow chart useful in explaining the proxy data disassembly processing.

In Step 161, the management protocol proxy data, and the address of the management protocol proxy as the other party which has transmitted the management protocol proxy data of interest are received from the interproxy communication part. In Step 162, the transmission source address, the transmission destination address and the SNMP PDU are fetched from the management protocol proxy data. In Step 163, the transmission source address and the transmission destination address which have been fetched from the management protocol proxy data are stored in the IP header part of the SNMP message. In Step 164, the SNMP PDU which has been fetched is stored in the SNMP message.

In a manner as described above, the proxy data assembly/disassembly unit disassembles the management protocol proxy data.

Fig. 19 is a block diagram showing a configuration of the address translation processing unit 63.

The address translation processing unit 63 includes: an SNMP message transmission source/transmission destination address translation unit 85 for translating the address of the transmission source and the address of the transmission destination of the SNMP message; and a PDU address translation unit 80 for translating the address information contained in the SNMP PDU. In addition, the PDU address translation

unit 80 includes: a PDU analysis/translation unit 81 for processing the analysis of the PDU and the address translation; an object identifier address translation unit 82 for processing the translation of the address which is contained as the object identifier of the address information contained in the PDU; an MIB instance value address translation unit 83 for processing the translation of the address which is contained as the MIB instance value of the address information contained in the PDU; and an agent-addr translation unit 84 for processing the translation of the address which is contained as the source address of trap (agent-addr) address of the address information contained in the PDU.

At the time when the management protocol proxy data has been delivered from the proxy data assembly/disassembly unit 62 to the address translation processing unit 63, first of all, the SNMP message transmission source/transmission destination address translation unit 85 carries out the address translation with respect to the transmission source and transmission destination of the SNMP message in the management protocol proxy data in accordance with the address translation definition 65. Next, the SNMP message transmission source/transmission destination address translation unit 85 delivers the management protocol proxy data to the PDU analysis/translation unit 81. Then, the PDU analysis/translation unit 81 carries out

the MIB instance value address translation unit 83. In addition, the object identifier of the MIB is extracted from the PDU to be delivered to the object identifier address translation unit 82. Then, the object

5 identifier address translation unit 82 subjects the IP address contained in the object identifier to the address translation in accordance with the ASN.1MIB define statement and the address translation definition 65 to return the resultant address information of the 10 object identifier back to the PDU analysis/translation unit 81. Then, the PDU analysis/translation unit 81 replaces the object identifier part of the MIB of the PDU with the address, after completion of the address translation, which has been received from the object 15 identifier address translation unit 82. Finally, the PDU analysis/translation unit returns the management protocol proxy data containing therein the PDU after completion of the address translation back to the proxy data assembly/disassembly unit 62.

20 In a manner as described above, the address translation processing unit can carry out the address translation with respect to the data of the management protocol.

Fig. 20 is a block diagram showing a configuration of the object identifier address translation 25 unit 82.

The object identifier address translation unit 82 includes: an ASN.1MIB define statement decoding

unit 88 for decoding the ASN.1MIB define statement 65;
a translation subject object identifier extraction unit
89 for extracting the object in which the object
identifier needs to be translated on the basis of the
5 definition contents of the decoded MIB; an object
identifier comparison unit 86 for comparing the object
identifier delivered from the PDU analysis/translation
unit 81 with the object identifier extracted by the
translation subject object identifier extraction unit
10 89 to judge whether or not the delivered object
identifier needs to be translated; and an address
translation extraction unit 87 for on the basis of the
definition information extracted by the translation
subject object identifier extraction unit 89 and the
15 address translation definition 65, subjecting the
object identifier to the address translation.

First of all, the ASN.1MIB define statement
decoding unit 88 reads out the ASN.1MIB define state-
ment 65 to decode the ASN.1MIB define statement 65 thus
20 read out to deliver the information of the MIB defini-
tion obtained by the decoding to the translation
subject object identifier extraction unit 89. Then,
the translation subject object identifier extraction
unit 89 extracts the MIB object having the possibility
25 of containing the IP address in the object identifier
from the delivered MIB statement to deliver the list of
the object identifiers of the corresponding MIB object
and the INDEX information as the definition information

of the corresponding MIB object to the object identifier comparison unit 86 and the address translation extraction unit 87, respectively. Now, by the MIB object having the possibility of containing the IP address in the object identifier is meant the MIB object representing the MIB table, i.e., the MIB object in which one or more MIB objects within the MIB table used as INDEX of the table are the IP addresses. Such an object is such that while when acquiring the instance as the value of the MIB object on the basis of the GET request or the like, INDEX as the instance identifier is added on the heels of the object identifier of the MIB object to be specified as the object identifier, since the IP address is used in INDEX at this time, there is a possibility that the IP address is contained in the object identifier. In addition, as for the INDEX information which is to be delivered to the address translation execution unit 87, in order that when a plurality of MIB objects are made correspond to one another as INDEXs of the MIB table, a part of the IP address of them may be translated, the information is delivered in which SYNTAXs as the kinds of MIB objects used as INDEXs are arranged in order. For example, in the case of the MIB table in which one MIB object of an integral number and the IP address are taken as INDEX, the instance identifier becomes the identifier having one sub-identifier for an integral number and four sub-identifiers for the IP address,

i.e., five sub-identifiers in total. In the address translation, since the second to fifth sub-identifiers of them need to be interpreted as the IP address to be translated, for INDEX, it is necessary to deliver the
5 information having a set of one integral number and one IP address to the address translation execution unit 87.

Now, at the time when the object identifier in the PDU has been delivered from the PDU analysis/
10 translation unit 81 to the object identifier address translation unit 82, first of all, the object identifier comparison unit 86 receives the object identifier. Then, the object identifier comparison unit 86 compares the object identifier delivered from
15 the PDU analysis/translation unit 81 with the object identifier list of the translation subject extracted by the translation subject object identifier extraction unit 89. If the object identifier delivered from the PDU analysis/translation unit 81 is contained in the
20 object identifier list of the translation subject, then the object identifier is delivered from the PDU analysis/translation unit 81 to the address translation execution unit 87. On the other hand, if the object identifier delivered from the PDU analysis/translation
25 unit 81 is not contained in the object identifier list of the translation subject, then the object identifier delivered from the PDU analysis/translation unit 81 is not translated at all to be returned back to the PDU

analysis/translation unit 81 as it is.

Next, with respect to the object identifier delivered therefrom, first of all, the address translation execution unit 87 specifies the location of the IP address appearing in the object identifier, i.e., the translation location on the basis of the INDEX information delivered from the translation subject object extraction unit 89 and next, carries out the address translation on the basis of the address translation definition 65 to return the object identifier after completion of the address translation back to the PDU analysis/translation unit 81.

In a manner as described above, the object identifier address translation unit can subject the IP address contained in the object identifier of the MIB to the address translation.

The processing in the PDU analysis/translation unit will hereinbelow be described with reference to a flow chart shown in Fig. 21.

In Step 111, the data representing the kind of PDU is extracted from the PDU. In Step 112, the address translation direction is determined from the kind of PDU. By the address translation direction is meant whether the address in the PDU is translated from the management address into the real address or from the real address into the management address. For the PDU which is transmitted from the managing side to the managed side, the management address is translated into

the real address. For the PDU transmitted from the managed side to the managing side, the real address is translated into the management address. Whether the PDU is transmitted from the managing side to the managed side or from the managed side to the managing side is determined every kind of PDU, and hence the address translation direction can be determined from the kind of PDU in accordance with a table shown in Fig. 22. In Step 113, it is judged whether or not the kind of PDU is the SNMP trap. If it is judged in Step 113 that the kind of PDU is the SNMP trap, then the processing proceeds to Step 114. On the other hand, if it is judged in Step 113 that the kind of PDU is not the SNMP trap, then the processing proceeds to Step 117. In Step 114, the trap transmission source address is extracted from the PDU. In Step 115, the trap transmission source address extracted in Step 114 and the information of the address translation direction determined in Step 112 are delivered to the trap transmission source address translation unit and then the trap transmission source address after completion of the translation is received. In Step 116, the trap transmission source address of the PDU is replaced with the trap transmission source address, after completion of the translation, which has been received in Step 115. In Step 117, it is judged whether or not variableBindingList is present in the PDU. If it is judged in Step 117 that variableBindingList is present

in the PDU, then the processing proceeds to Step 118.

On the other hand, if it is judged in Step 117 that variableBindingList is absent in the PDU, then the processing in the PDU analysis/translation unit is

5 completed. In Step 118, one unprocessed variableBind is executed from variableBindingList. In Step 119, the object identifier and the value are extracted from variableBind extracted in Step 118. In Step 120, both of the MIB instance value extracted in Step 119 and the
10 information of the translation direction determined in Step 112 are delivered to the MIB instance value address translation unit and then the MIB instance value after completion of the translation is received.

In Step 121, both of the object identifier extracted in
15 Step 119 and the information of the translation direction determined in Step 112 are delivered to the object identifier address translation unit and then the object identifier after completion of the translation is received. In Step 122, the object identifier of

20 variableBind of the PDU is replaced with the object identifier, after completion of the translation, which has been received in Step 121, and also the MIB instance value of variableBind is replaced with the MIB instance value, after completion of the translation,

25 which has been received in Step 120. In Step 123, it is judged whether or not the unprocessed variable-Binding still remains. If it is judged in Step 123 that the unprocessed variableBinding still remains,

then the processing proceeds to Step 118. On the other hand, if it is judged that the unprocessed variable-Binding does not yet remain, then the processing in the PDU analysis/translation unit is completed.

5 In a manner as described above, the address information in the PDU can be translated. Next, the processing in the translation subject object identifier extraction unit will hereinbelow be described with reference to a flow chart shown in Fig. 23.

10 In Step 131, one object identifier which is defined in the MIB statement is fetched. In Step 132, it is judged whether or not the object identifier fetched in Step 131 is the identifier representing the MIB table. If it is judged in Step 132 that the object
15 identifier fetched in Step 131 is the identifier representing the MIB table, then the processing proceeds to Step 133. On the other hand, if it is judged in Step 132 that the object identifier fetched in Step 131 is not the identifier representing the MIB
20 table, then the processing proceeds to Step 136. In Step 133, SYNTAX of the MIB object in the table which is specified as INDEX of the object identifier is extracted. In Step 134, it is judged whether or not there is even one object in which SYNTAX of the MIB
25 object in the table extracted in Step 133 is SYNTAX representing the IP address. If it is judged in Step 134 that there is even one object in which SYNTAX is SYNTAX representing the IP address, then the processing

proceeds to Step 135. On the other hand, if it is judged in Step 134 that there is not even one object in which SYNTAX is SYNTAX representing the IP address, then the processing proceeds to Step 136. In Step 135,
5 it is judged that the object identifier fetched in Step 131 is the object identifier of a subject of the translation.

On the other hand, in Step 136, it is judged that the object identifier fetched in Step 131 is not
10 the object identifier of a subject of the translation. In Step 137, it is judged whether or not any of the unprocessed object identifiers still remains in the MIB define statement. If it is judged in Step 137 that any of the unprocessed object identifiers still remains in
15 the MIB define statement, then the processing proceeds to Step 131. On the other hand, if it is judged in Step 137 that any of the unprocessed object identifiers does not yet remains in the MIB define statement, then the processing proceeds to Step 138. In Step 138, the
20 object identifier comparison unit is informed of all of the object identifiers each of which has been judged to be a subject of the translation. In Step 139, the address translation execution unit is informed of all of the INDEX information of the object identifiers each
25 of which has been judged to be a subject of the translation.

In a manner as described above, it is possible to realize the processing in the translation

subject object identifier extraction unit.

Fig. 24 shows a definition example of the address translation definition 65.

A definition line 301 is a definition example
5 in the case where only a first octet of the IP address
is translated. In this case, with respect to all of
the real addresses in each of which the first octet is
10, the real address in which the first octet is
translated into 100 becomes the management address.
10 For example, in the case where the address 100.1.2.3 is
present in the data of the management protocol which
has been transmitted from the manager 40, the address
of interest is translated into an address 10.1.2.3
through the address translation to be relayed to the
15 managed node 50. On the contrary, in the case where
the address 10.1.2.3 is present in the response from
the managed node 50, the address of interest is
translated into an address 100.1.2.3 through the
address translation to be relayed to the manager 40. A
20 definition line 302 is a definition example in the case
where a first octet and a second octet of the IP
address are both translated. In this case, with
respect to all of the real addresses in each of which
the first octet is 172 and the second octet is 16, the
25 real address in which the first octet is translated
into 101 and the second octet is translated into 10
becomes the management address. For example, in the
case where the address 101.10.1.2 is present in the

data of the management protocol which has been transmitted from the manager 40, the address of interest is translated into an address 172.16.1.2 through the address translation to be relayed to the managed node 50. On the contrary, in the case where the address 172.16.1.2 is present in the response from the managed node 50, the address of interest is translated into an address 101.10.1.2 through the address translation to be relayed to the manager 40. A definition line 303 is a definition example in the case where a first octet, a second octet and a third octet of the IP address are all translated. In this case, with respect to all of the real addresses in each of which the first octet is 172, the second octet is 17 and the third octet is 50, the real address in which the first octet is translated into 110, the second octet is translated into 20 and the third octet is translated into 80 becomes the management address. For example, in the case where the address 110.20.80.1 is present in the data of the management protocol which has been transmitted from the manager 40, the address of interest is translated into an address 172.17.50.1 through the address translation to be relayed to the managed node 50. On the contrary, in the case where the address 172.17.50.1 is present in the response from the managed node 50, the address of interest is translated into an address 110.20.80.1 through the address translation to be relayed to the manager 40.

In addition, a definition line 304 is a definition example in the case where when all of the octets from the first octet to the fourth octet are translated. In this case, the real address in which the address is
5 192.168.20.5 is translated into the address 120.60.11.8 for management. For example, in the case where the address 120.60.11.8 is present in the data of the management protocol which has been transmitted from the manager 40, the address of interest is translated into
10 an address 192.168.20.5 through the address translation to be relayed to the managed node 50. On the contrary, in the case where the address 192.168.20.5 is present in the response from the managed node 50, the address of interest is translated into an address 120.60.11.8
15 through the address translation to be relayed to the manager 40.

Fig. 25 is a diagram showing an example of a configuration of another virtual network management system.

20 In this case, while the manager 40 and the management protocol proxy 60a on the manager side are present on the global network 10, it is possible to carry out the virtual network management using the management address without being different from the
25 case having the configuration shown in Fig. 1.

Fig. 26 and Fig. 27 are respectively block diagrams showing configurations of the management protocol proxies in still another embodiment and yet

another embodiment, i.e., embodiments in the case where the address translation is carried out by the management protocol proxy on the managing side.

In these cases as well, the address translation can be carried out in the same manner as that in the case where the address translation is carried out by the management protocol proxy on the managed side shown in Fig. 12 and Fig. 13. But, when the address translation is carried out in the management protocol proxy on the managing side, the address translation definition which is to be carried out in the management protocol proxy on the managing side is defined individually every private network on the managed side. For example, the address translation definition 65b and the address translation definition 65c are respectively defined for the private network B on the managed side and for the private network C on the managed side. As a result, even in the case where the private addresses of the managers conflict with each other in a plurality of private networks on the managed side as in the configuration shown in Fig. 9, the address translation can be made carried out properly.

Fig. 28 is a block diagram showing a configuration of the virtual network management system in a further embodiment, and Fig. 29 is a block diagram showing a configuration of the management protocol proxy on the managed side in the further embodiment. In these figures, each of the configurations is such

that the management protocol proxy on the managed side is operated on the same apparatus as the NAT conforming to RFC1631.

The configuration of this example is such
5 that the manager 40 and the management protocol proxy 60a on the managing side are present on the global network 10, and the management address translation is carried out in the management protocol proxy on the managed side, and the management protocol proxy on the
10 managed side is operated on the same apparatus as that of the NAT 20 conforming to RFC 1631.

Fig. 29 shows the configuration of the management protocol proxy 60b on the managed side which is operated on the same apparatus as that of the NAT 20
15 not having the address translation function of the management protocol conforming to RFC1631. A point of difference from the configuration shown in Fig. 12 is that when the interproxy communication unit 61 is intended to communicate with the management protocol
20 proxy 60a on the managing side, the communication is carried out through a global network side communication protocol processing unit 21, and when the SNMP message transmission/reception unit 66 transmits/receives the data to/from the managed node 50, the communication is
25 carried out through a private network side communication protocol processing unit 23 of the NAT 20. By the way, as for the motion conforming to RFC1631 of the NAT 20, with respect to the packet which is intended to

pass through the associated apparatuses of the NAT 20 from the global network side to the private network side, first of all, the global network side communication protocol processing unit 21 catches the packet
5 which is intended to pass therethrough to deliver the packet thus caught to an RFC1631 conformity address translation processing unit 22, and then the RFC1631 conformity address translation processing unit 22 carries out the address translation. Thereafter, the
10 packet of interest is delivered to the private network side communication protocol processing unit 23 to be sent to the private network side by the private network side communication protocol processing unit 23.

On the contrary, with respect to the packet
15 which is intended to pass through the associated apparatuses of the NAT 20 from the private network side to the global network side, first of all, the private network side communication protocol processing unit 23 catches the packet which is intended to pass there-
20 through to deliver the packet thus caught to the RFC1631 conformity address translation processing unit 22 and then the RFC1631 conformity address translation processing unit 22 carries out the address translation. Thereafter, the packet is delivered to the global
25 network side communication protocol processing unit 21 to be sent to the global network side by the the global network side communication protocol processing unit 21.

However, the communication made by the

interproxy communication unit 61 is the communication in which the address on the global network side of the apparatus in which the NAT and the management protocol proxy on the managed side are both operated is made
5 either the transmission destination or the transmission source, but is not the communication in which the associated information is not intended to pass through the NAT 20. For this reason, the data is delivered from the global network side communication processing
10 unit 21 to the interproxy communication unit 61 as it is without passing through the RFC1631 conformity address translation processing unit 22.

In addition, with respect to the communication as well made by the SNMP message transmission/
15 reception unit 66, it is the communication in which the address on the private network side of the apparatus in which the NAT and the management protocol proxy on the managed side are both operated is made either the transmission destination or the transmission source,
20 but is not the communication in which the associated information is not intended to pass through the NAT 20. For this reason, the data is delivered from the private network side communication processing unit 23 to the SNMP message transmission/reception unit 66 as it is
25 without passing through the RFC1631 conformity address translation processing unit 22.

From the foregoing, it is possible that the management protocol proxy having the same configuration

as that of the management protocol proxy on the managed side shown in Fig. 12 is employed to be operated on the same apparatus as the NAT 20 to realize the virtual network management.

5 By the way, as has already been described, the processing in the management protocol proxy, as shown in the flow chart, can be realized in the form of the program.

 Above, while the NAT has been described as
10 the apparatus for carrying out the address translation of the IP layer (the NAT does not have the function of translating the address of the PDU part of the management protocol), in the case where the NAT has the function of carrying out the address translation of the
15 IP layer and the function of translating the address of the PDU part of the management protocol, the NAT and the management protocol proxy server can be selectively utilized to carry out the address translation of the PDU part of the management protocol.

20 As set forth hereinabove, according to the present invention, the data communication can be carried out between the proxies of the management protocol, and also the data communication by the management protocol can be carried out between the
25 nodes each not having the global address.

 Although the present invention has been described with reference to specific embodiments, these descriptions are not meant to be construed in a

limiting sense. Various modifications of the disclosed
embodiments, as well as alternative embodiments of the
invention will become apparent to persons skilled in
the art upon reference to the description of the inven-
5 tion. It should be appreciated by those skilled in the
art that the conception and the specific embodiment
disclosed may be readily utilized as a basis for
modifying or designing other configurations for carry-
ing out the same purpose of the present invention. It
10 should also be realized by those skilled in the art
that such equivalent configurations do not depart from
the spirit and scope of the invention as set forth in
the appended claims.

It is therefore contemplated that the claims
15 will cover any such modifications or embodiments that
fall within the true scope of the invention.